

Copredication in Homotopy Type Theory

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Introduction

Copredication:

the phenomenon where two or more predicates with different requirements on their arguments are applied to a single argument

- a) The lunch was delicious but took forever. (Asher, 2011)
- b) The heavy book is easy to understand. (Gotham, 2014)
- c) John picked up and mastered three books. (Asher 2011)

History

- Asher(2011), nouns as "dot objects" (Pustejovsky, 1995)
 - i.e., objects that can be viewed under different "aspects"
- Cooper (2011)
 - nouns as functions from "records" to "record types"
- Luo (2012), Chatzikyriakidis and Luo (2012, 2013, 2015)
 - common nouns as types
 - dot types + coercive subtyping
- Gotham(2014)
 - `book' denotes the set of composite objects physical+informational
 - criteria of individuation are combined during semantic composition

Montague semantics

(1) a. The lunch was delicious.

b. The lunch took forever.

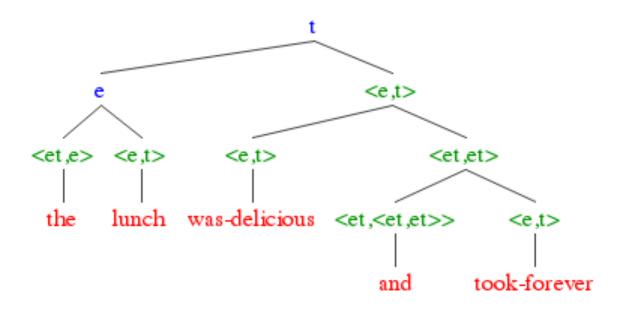
c. The lunch was delicious and took forever.

$$\begin{bmatrix} \text{was delicious} \end{bmatrix} = \lambda x \in D_e.x \text{ was delicious} \\ \begin{bmatrix} \text{took forever} \end{bmatrix} = \lambda x \in D_e.x \text{ took forever} \\ \begin{bmatrix} \text{and} \end{bmatrix} = \begin{bmatrix} \lambda f \in D_{\langle e,t \rangle}. \\ [\lambda g \in D_{\langle e,t \rangle}.[\lambda x \in D_e.f(x) = g(x) = 1]] \end{bmatrix} \\ \begin{bmatrix} \text{lunch} \end{bmatrix} = \lambda x \in D_e.x \text{ is a lunch} \\ \\ \begin{bmatrix} \text{the} \end{bmatrix} = \lambda f \in D_{\langle e,t \rangle} \& \exists ! x \in D_e[f(x) = 1]. \\ !y[f(y) = 1], \\ \text{where } \exists ! x[f(x) = 1] \\ abbreviates "there is exactly \\ \text{one x such that } f(x)=1 \text{ and } !y[\phi] \text{ returns} \\ "that unique y such that } f(y)=1". \end{bmatrix}$$

Functional Application (FA): if α is a branching node with β and γ as its daughters, then α is in the domain of $\llbracket \cdot \rrbracket$ if both β and γ are, and if $\llbracket \gamma \rrbracket$ is in the domain of $\llbracket \beta \rrbracket$. In this case, $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket (\llbracket \gamma \rrbracket)$ (Heim and Kratzer[1998]).

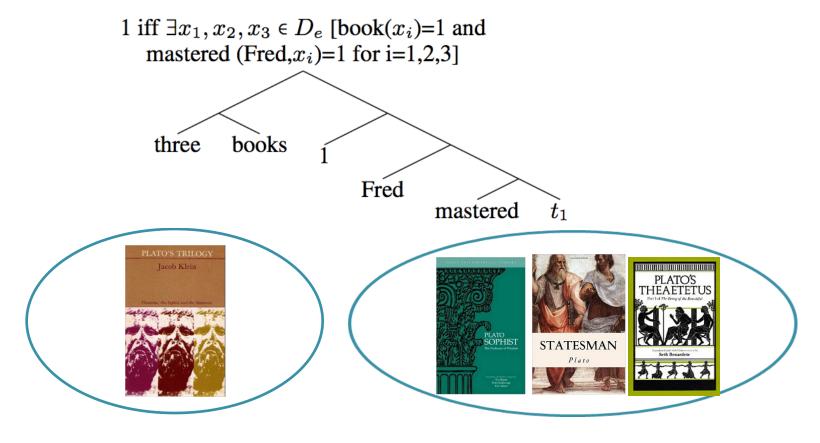


- (1) a. The lunch was delicious.
 - b. The lunch took forever.
 - c. The lunch was delicious and took forever.



Montague semantics

- (2) a. Fred picked up three books.
 - b. Fred mastered three books.
 - c. Fred picked up and mastered three books.



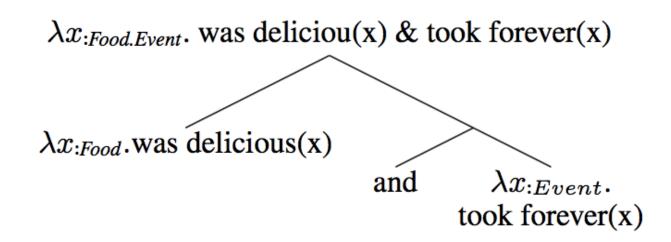
- (1) a. The lunch was delicious.
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A.B is only well-formed if A and B do not share common components, and both projections, one from A.B to A and the other from A.B to B, are coercions in the coercive subtyping framework.(Luo[2012])

- a. Food.Event $<_c$ Food
- b. Food.Event $<_c$ Event
- c. Lunch $<_c$ Food.Event $<_c$ Food
- d. Lunch $<_c$ Food.Event $<_c$ Event

- (1) a. The lunch was delicious.
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 - a. Food.Event $<_c$ Food
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 - c. Lunch $<_c$ Food.Event $<_c$ Food
 - d. Lunch $<_c$ Food.Event $<_c$ Event
 - a. Food \rightarrow Prop \leq_c Food.Event \rightarrow Prop \leq_c Lunch \rightarrow Prop
 - b. Event \rightarrow Prop \leq_c Food.Event \rightarrow Prop \leq_c Lunch \rightarrow Prop

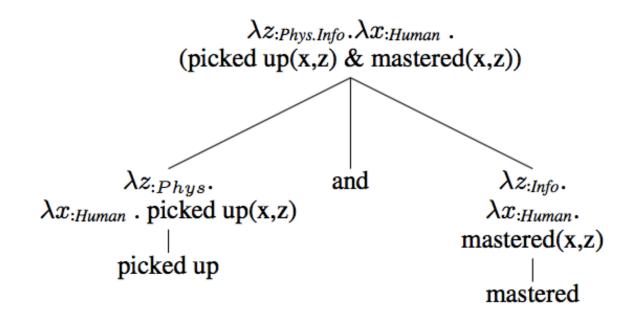
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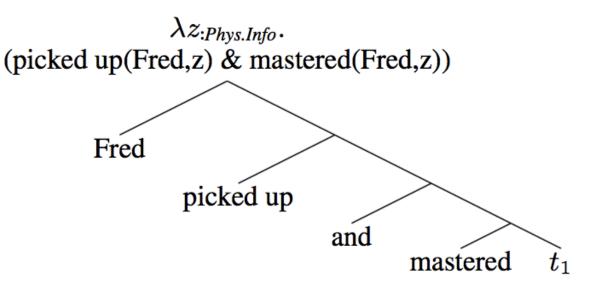
- (2) a. Fred picked up three books.
 - b. Fred mastered three books.
 - c. Fred picked up and mastered three books.

 $\lambda B_{:Book \rightarrow Prop} \exists x, y, z_{:Book} [B(x)B(y)B(z)]$ $\lambda A_{:Type} \cdot \lambda B_{:A \rightarrow Prop} \cdot \exists x, y, z_{:A}$ [B(x)B(y)B(z)] $Book_{:Type}$ books three

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 $\exists x, y, z_{:Book}$ [picked up(Fred,x) and mastered(Fred,x) picked up(Fred,y) and mastered (Fred,y) picked up(Fred,z) and mastered (Fred,z)]

 $\begin{array}{c} \lambda B_{:Book \rightarrow Prop}. \\ \exists x, y, z_{:Book} \ [B(x) \ B(y) \ B(z)] \\ \downarrow \\ \text{three books} \end{array}$

 $\lambda z_{:Phys.Info}$. (picked up(Fred,z) and mastered(Fred,z)) Fred picked up and mastered t_1

a. Book $<_{c_1}$ Physical

b. Book $<_{c_2}$ Informational

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 $\exists x, y, z_{:Book}$ [picked up(Fred,x) & mastered(Fred,x) picked up(Fred,y) & mastered (Fred,y) picked up(Fred,z) & mastered (Fred,z)]

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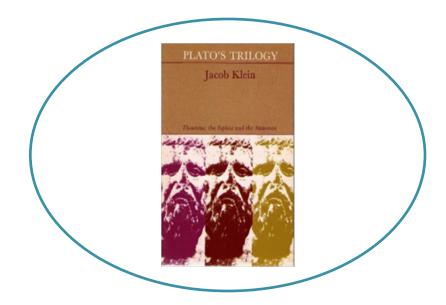
 $\exists x, y, z_{:Book}$ [picked up(Fred,x) & mastered(Fred,x) picked up(Fred,y) & mastered (Fred,y) picked up(Fred,z) & mastered (Fred,z)]

Variable PHY:forall x:Book, forall y:Book, not(x=y:>Book)-> not(x=y:>Phy). Variable INFO:forall x:Book, forall y:Book, not(x=y:>Book)-> not(x=y:>Info).



The problem

Fred mastered three books.





The problem

Five Books are heavy but easy to understand.(Gotham[2012])



Homotopy Type Theory

- *intensional* version of Martin-Löf's type theory [ML75]
- a *proof-relevant* interpretation of equality
- *propositions-as-types* principle
- a full cumulative hierarchy of universes
- judgmental vs propositional equality

Homotopy Type Theory

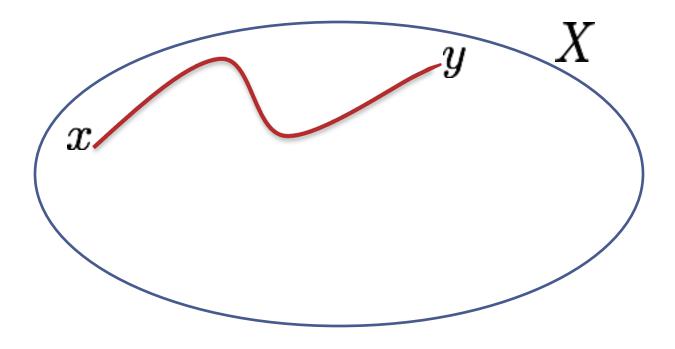
- \bullet variables x, x', . . .
- primitive constants c, c', . . .
- defined constants f , f', . . .
- $t \coloneqq x | \lambda x.t | t(t') | c | f$

some primitive constants:

- a hierarchy of universes $\mathcal{U}_1, \mathcal{U}_2, \dots$
- dependent function types $\Pi_{a:A}B$
- dependent pair types $\Sigma_{a:A}B$
- identity types $a =_A b, A =_U B$

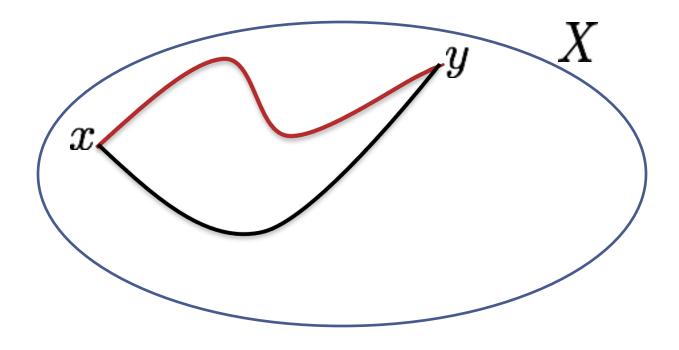


$p: [0,1] \to X$ where p(0) = x and p(1) = y



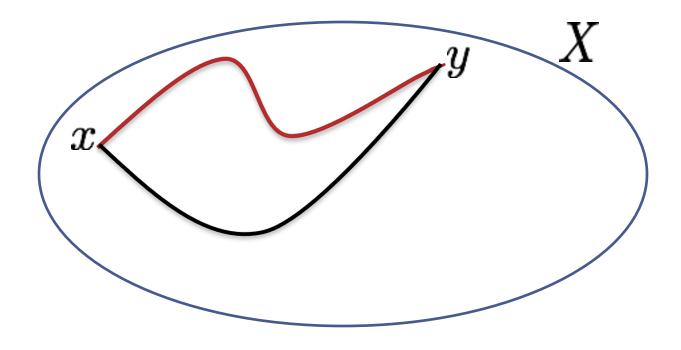


$p: [0,1] \to X$ where p(0) = x and p(1) = y





 $x =_X y$





A = B



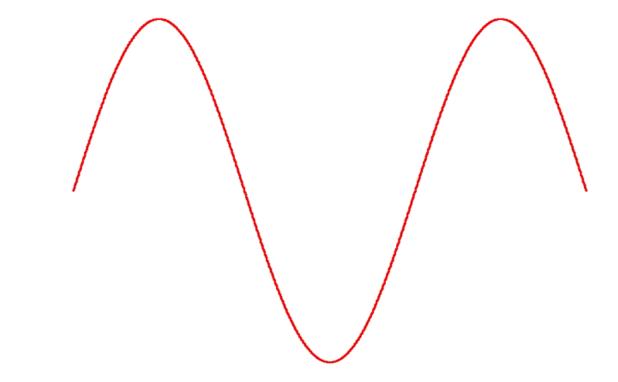
=?



A homotopy between a pair of continuous maps $f: X1 \to X2$ and $g: X1 \to X2$ is a continuous map $H: X1 \times [0,1] \to X2$ satisfying H(x,0) = f(x) and H(x,1) = g(x).

If there is such a function H then we say f and g are homotopic, written $f \sim g$.

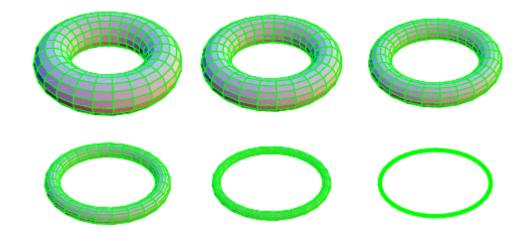




Johnson, Chris. "Path Homotopy Animation." Youtube. Youtube, 11 May 2009. Web. 9 July 2016.

Two spaces X and Y are homotopy equivalent if there are maps $f: X \to Y$ and $f': Y \to X$ such that $f' \circ f \sim idX$ and $f \circ f' \sim idY$.





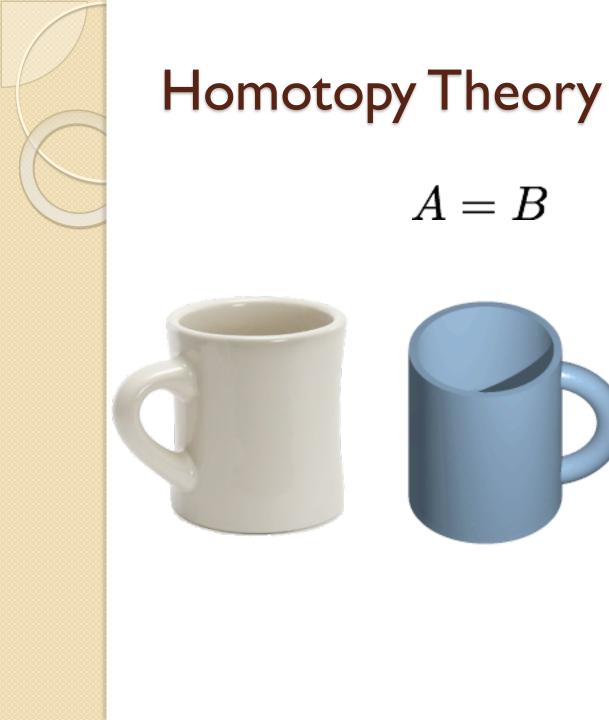
Rowland, Todd. "Homotopic." From MathWorld--A Wolfram Web Resource, created by Eric W.Weisstein. http://mathworld.wolfram.com/Homotopic.html

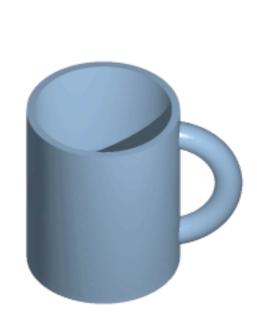


A = B





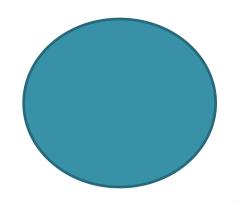




A = B

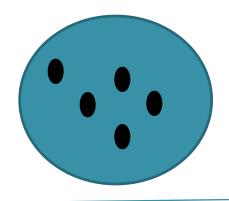


















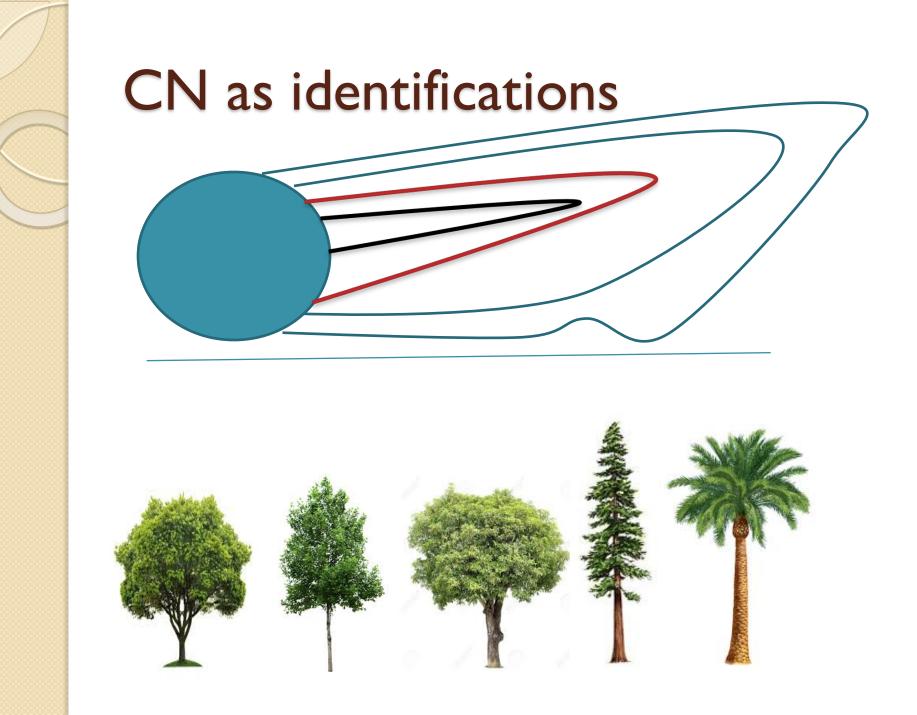


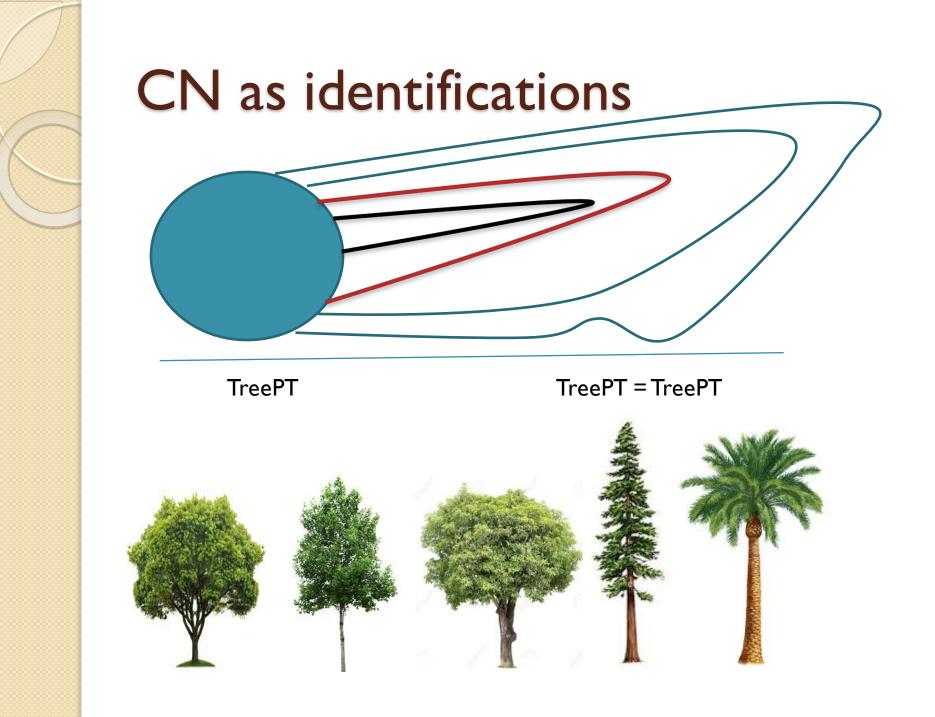
Sameness

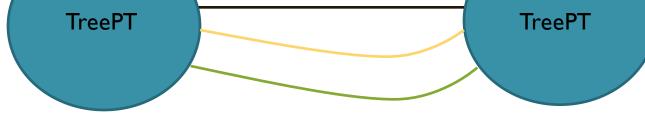








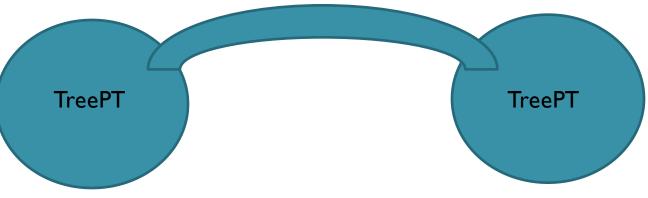




TreePT = TreePT







TreePT = TreePT





when X has no aspects, X is :

$$X_{PT} = X_{PT}$$



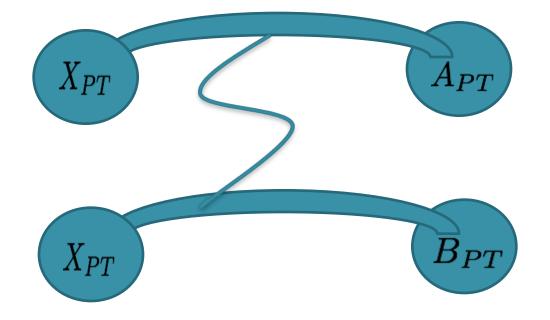
when X has one aspect, namely A, X is:

 $X_{PT} = A_{PT}$

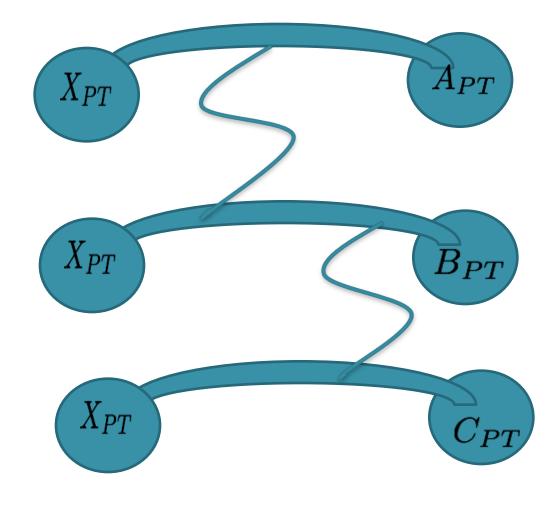


when X has two aspects, namely A and B, X is:

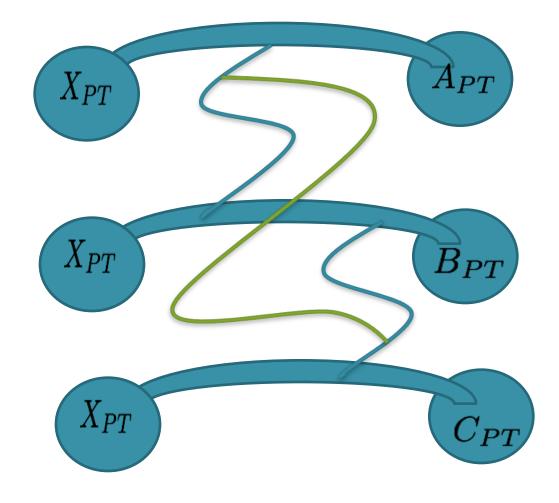
$$(X_{PT} = A_{PT}) = (X_{PT} = B_{PT})$$

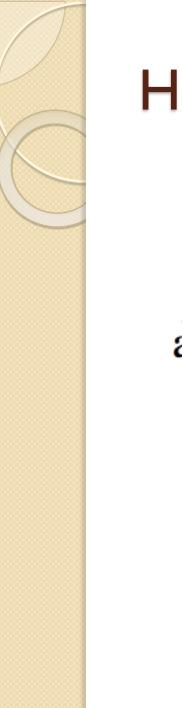


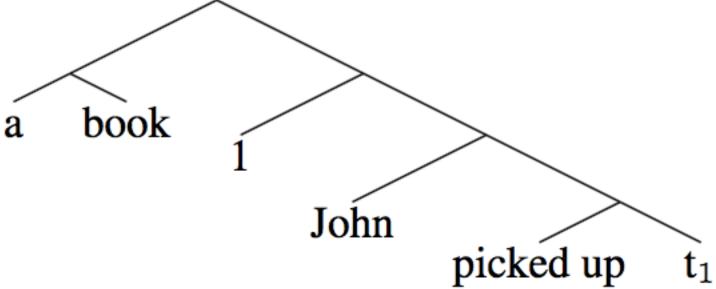
when X has three aspects, namely A, B and C, X is:



when X has three aspects, namely A, B and C, X is:







- = $[a \operatorname{book}]^g (\lambda b. [John picked up t_1]^{g[1 \to b]})$ PA
- = $[a \operatorname{book}]^g (\lambda b. [\operatorname{picked} \operatorname{up} t_1]^{g[1 \to b]} ([John]^{g[1 \to b]}))$ FA
- $= \llbracket a \operatorname{book} \rrbracket^{g} (\lambda b. \llbracket \operatorname{picked} \operatorname{up} t_1 \rrbracket^{g[1 \to b]} (John))$ = $\llbracket a \operatorname{book} \rrbracket^{g} (\lambda b. \llbracket \operatorname{picked} \operatorname{up} \rrbracket^{g[1 \to b]} (\llbracket t_1 \rrbracket^{g[1 \to b]}) (John))$ Lex John
- FA
- = $[a \operatorname{book}]^g (\lambda b. [picked up]]^{g[1 \to b]}(b) (John))$

- $= \begin{bmatrix} a \text{ book} \end{bmatrix}^{g} (\lambda b. \begin{bmatrix} \text{John picked up } t_1 \end{bmatrix}^{g[1 \to b]})$ PA $= \begin{bmatrix} a \text{ book} \end{bmatrix}^{g} (\lambda b. \begin{bmatrix} \text{picked up } t_1 \end{bmatrix}^{g[1 \to b]} (\begin{bmatrix} \text{John} \end{bmatrix}^{g[1 \to b]}))$ FA $= \begin{bmatrix} a \text{ book} \end{bmatrix}^{g} (\lambda b. \begin{bmatrix} \text{picked up } t_1 \end{bmatrix}^{g[1 \to b]} (John))$ Lex John $= \begin{bmatrix} a \text{ book} \end{bmatrix}^{g} (\lambda b. \begin{bmatrix} \text{picked up } \end{bmatrix}^{g[1 \to b]} (\begin{bmatrix} t_1 \end{bmatrix}^{g[1 \to b]}) (John))$ FA $= \begin{bmatrix} a \text{ book} \end{bmatrix}^{g} (\lambda b. \begin{bmatrix} \text{picked up } \end{bmatrix}^{g[1 \to b]} (\begin{bmatrix} t_1 \end{bmatrix}^{g[1 \to b]}) (John))$ FA
- = $[a \operatorname{book}]^g (\lambda b. [picked up]]^{g[1 \to b]}(b) (John))$

Lemma 1. Given an element b of type Book, the physical component of b exists and it is an inhabitant of type Physical.

Lemma 1. Given an element b of type Book, the physical component of b exists and it is an inhabitant of type Physical. Proof. Let b is an element of Book then b is a path from:

 $(BookPrototype =_{U_i} PhysicalPrototype)$

to (*BookPrototype* $=_{\mathcal{U}_i}$ *InformationalPrototype*). Every path from a type A to a type B induces two functions, one from A to B and the other from B to A. So the path b induces a function f_b from (BookPrototype = U_i InformationalPrototype) to (BookPrototype $=_{\mathcal{U}_i} PhysicalPrototype$). By ANN (BookPrototype $=_{\mathcal{U}_i}$ InformationalPrototype) and (BookPrototype) $=_{\mathcal{U}_i}$ PhysicalPrototype) are not empty and ANN gives us one member of each namely x and y, which are dependent on b. Then $f_b(x)$ is a path from BookPrototype to PhysicalPrototype. But $(f_b(x))^{-1}$ is a path from *PhysicalPrototype* to *BookPrototype*. The path $y \circ (f_b(x))^{-1}$ is a path from *PhysicalPrototype* to itself which means it is an element of the *Physical* type. $y \circ (f_b(x))^{-1}$ is an inhabitant of type Physical and it is determined by b therefore it can be considered as the physical component of b.

Lemma 1. Given an element b of type Book, the physical component of b exists and it is an inhabitant of type Physical.

book-is-physical : $(b : Book) \rightarrow Physical$ book-is-physical b =PhysicalPrototype = $\langle inverse-p(coe! (fst b) ((not-empty.bPoint (snd b)))) \rangle$ BookPrototype = $\langle (not-empty.aPoint (snd b)) \rangle$ PhysicalPrototype =

- = $[a \operatorname{book}]^g (\lambda b. [John picked up t_1]^{g[1 \rightarrow b]})$ PA $= [a \operatorname{book}]^{g} (\lambda b. [\operatorname{picked} \operatorname{up} t_1]^{g[1 \to b]} ([John]^{g[1 \to b]}))$
- FA Lex John
- $= \llbracket a \operatorname{book} \rrbracket^{g} (\lambda b. \llbracket \operatorname{picked} \operatorname{up} t_1 \rrbracket^{g[1 \to b]} (John))$ = $\llbracket a \operatorname{book} \rrbracket^{g} (\lambda b. \llbracket \operatorname{picked} \operatorname{up} \rrbracket^{g[1 \to b]} (\llbracket t_1 \rrbracket^{g[1 \to b]}) (John))$ FA
- = $[a \operatorname{book}]^g (\lambda b. [picked up]]^{g[1 \to b]}(b) (John))$

Functional Application(revised): if α is a constituent with β and γ as its daughters, then if $[\beta]$ is a function whose domain contains $[\gamma]$ or an aspect of $[\gamma]$, then $[\alpha] = [\beta]([\gamma]_{\beta})$

where $[\![\gamma]\!]_{\beta} = [\![\gamma]\!]$ if $[\![\gamma]\!]$ is in the domain of $[\![\beta]\!]$, and if $[\![\beta]\!]$ is a function whose domain contains an aspect $[\chi]$ of $[\gamma]$, $[\gamma]_{\beta}$ is the aspect $[\![\chi]\!]$.

- $= \llbracket a \text{ book} \rrbracket^g (\lambda b. \llbracket \text{John picked up } t_1 \rrbracket^{g[1 \to b]})$ PA
- $= \llbracket a \operatorname{book} \rrbracket^g (\lambda b. \llbracket \operatorname{picked} \operatorname{up} t_1 \rrbracket^{g[1 \to b]} (\llbracket John \rrbracket^{g[1 \to b]}))$ FA
- $= [a \operatorname{book}]^g (\lambda b. [\operatorname{picked} \operatorname{up} t_1]^{g[1 \to b]} (John)) \qquad \text{Lex John}$
- $= \llbracket a \operatorname{book} \rrbracket^{g} (\lambda b. \llbracket \operatorname{picked} \operatorname{up} \rrbracket^{g\llbracket 1 \to b}] (\llbracket t_1 \rrbracket^{g\llbracket 1 \to b}] (John)) \qquad \text{FA}$
- = $[a \operatorname{book}]^g (\lambda b. [picked up]]^{g[1 \to b]}(b) (John))$
- = $[\text{three book}]^g (\lambda b. [\text{picked up }]^{g[1 \to b]}(\varphi_b)(John))$

FA

 $= \llbracket a \operatorname{book} \rrbracket^g (\lambda b. [\lambda h.1 \operatorname{iff} h \operatorname{picked} \operatorname{up} \varphi_b] (John)) \quad \text{Lex picked} \\ \operatorname{up, } \lambda \text{-Conv.} \\ = \llbracket a \operatorname{book} \rrbracket^g (\lambda b.1 \operatorname{iff} \operatorname{John} \operatorname{picked} \operatorname{up} \varphi_b) \qquad \lambda \text{-Conv.} \\ = 1 \operatorname{iff} \operatorname{there exists} x \in Book \\ \operatorname{such that} \operatorname{John} \operatorname{picked} \operatorname{up} \varphi_x \qquad \operatorname{by} (20) \lambda \text{-Conv.} \end{cases}$

- (2) a. John picked up three books.
 - b. John mastered three books.
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[[three books]] = [[three]] ([[books]]) = $\lambda \gamma \in D_{\langle Book, U \rangle}$.1 iff there exist $x, y, z \in Book$ such that $\gamma(x), \gamma(y)$ and $\gamma(z)$ are inhabitable

- (2) a. John picked up three books.
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 $\llbracket (2a) \rrbracket^g = 1 \text{ iff there exist } x, y, z \in Book \text{ such that John picked up} \\ \varphi_x, \varphi_y \text{ and } \varphi_z \\ \llbracket (2b) \rrbracket^g = 1 \text{ iff there exist } x, y, z \in Book \text{ such that John mastered} \\ \rho_x, \rho_y \text{ and } \rho_z \\ \llbracket (2c) \rrbracket^g = 1 \text{ iff there exist } x, y, z \in Book \text{ such that John mastered} \\ \rho_x, \rho_y \text{ and } \rho_z \\ \llbracket (2c) \rrbracket^g = 1 \text{ iff there exist } x, y, z \in Book \text{ such that John mastered} \\ \rho_x, \rho_y \text{ and } \rho_z \text{ and John picked up } \varphi_x, \varphi_y \text{ and } \varphi_z \end{aligned}$

Semantic computation using HoTT logic

Theorem 4. $(h : Human)(a : \forall (P : Book \rightarrow U_i) \rightarrow (\Sigma Book P))$ $(p : Physical \rightarrow Human \rightarrow U_j) \rightarrow \Sigma Physical (\lambda x \rightarrow p x h)$

john-picked-up-a-book : (j : Human) $(a: \forall (P: \mathsf{Book} \to \mathsf{Type}))$ $\rightarrow (\Sigma \operatorname{Book} P))$ $(p: Physical \rightarrow Human \rightarrow Type _)$ $\rightarrow \Sigma$ Physical ($\lambda x \rightarrow p x j$) john-picked-up-a-book j a p =(a-book-is-a-physical a) (p-fix-middle j p)



Future work

- Formal semantics of natural languages
 "possible world"?
- HoTT

 subtyping?
- Homotopy
 - Prefiguration as theorem?

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